

University of Stuttgart Institute of Information Security

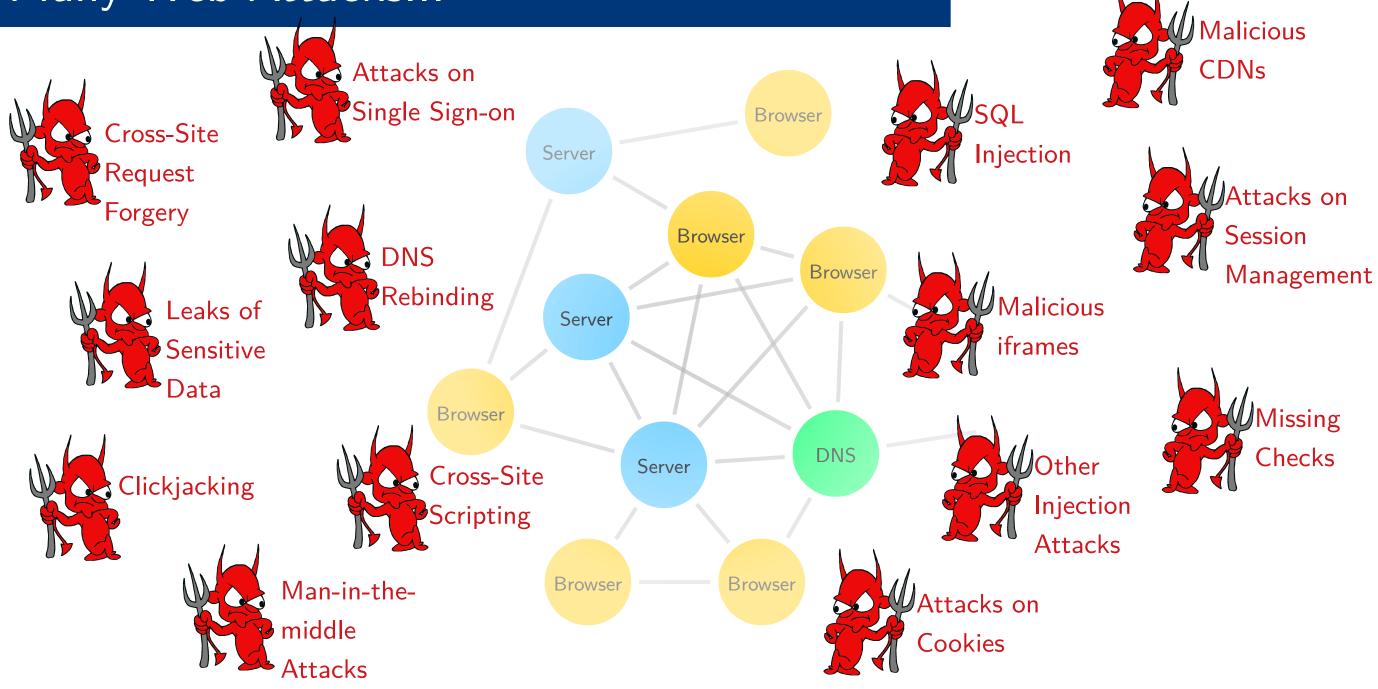
> WIM: An Expressive Formal Model of the Web Infrastructure

Ralf Küsters

2019-06-27

Joint work with Daniel Fett, Pedram Hosseyni, and Guido Schmitz [S&P 2014, ESORICS 2015, CCS 2015, CCS 2016, CSF 2017, S&P 2019]

Many Web Attacks...

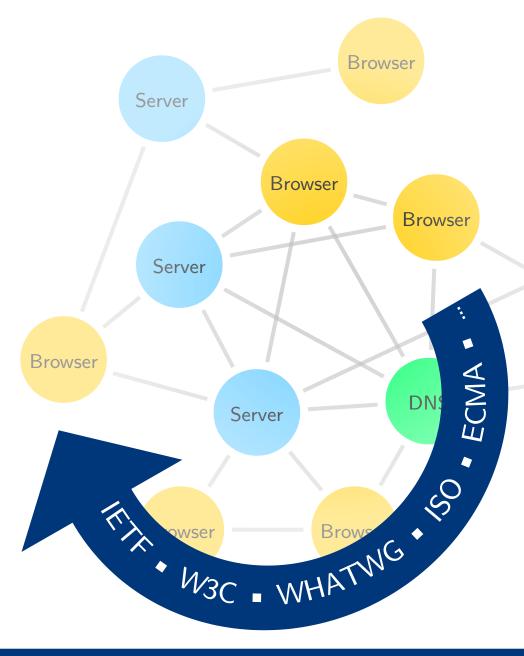


...but why?

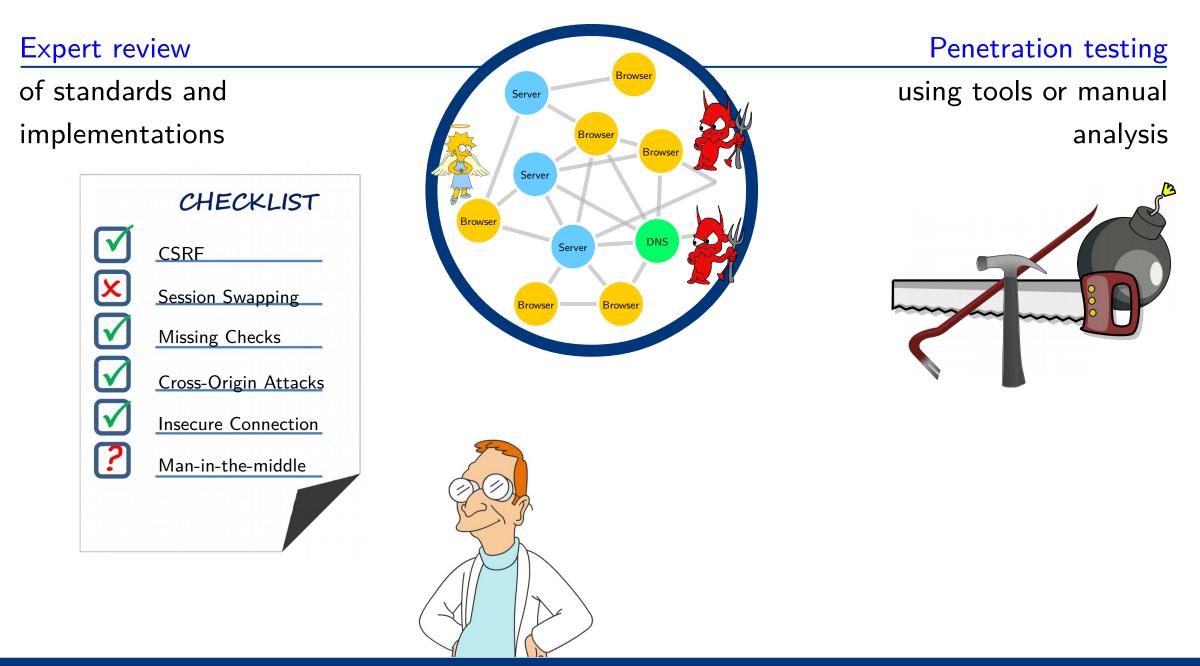
The web is complex ...

- Network of heterogeneous components
- Large number of complex standards developed at a high pace by many separate organizations

- ... and web applications as well.
- More features, more interaction
- Many bugs and errors



Finding Vulnerabilities: Current Practice

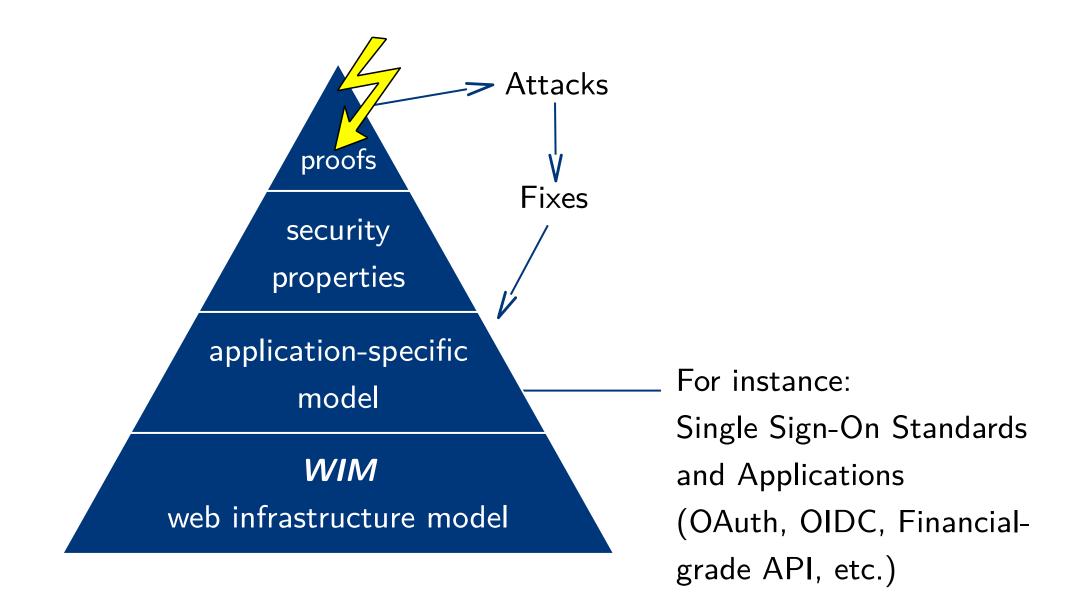


Downsides

- ► It is easy to miss attacks, even for experts
- Pentesting focuses on known attacks
- Finding new attack types depends on the creativity of the experts
- Both methods do not guarantee security, not even for a limited set of attacks

Can we develop a more systematic way of finding vulnerabilities, and even prove security (in a meaningful model of the web infrastructure as a whole)?

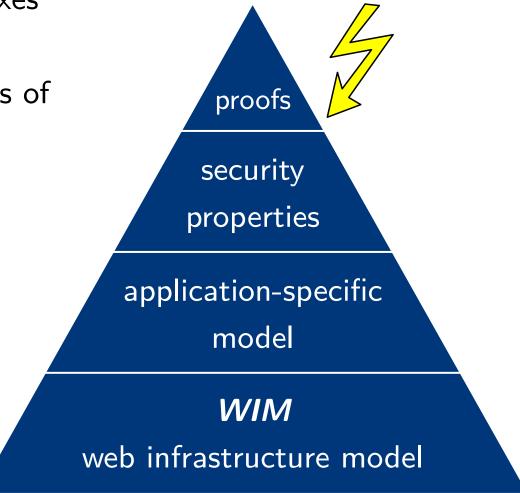
Our Model-Based Approach



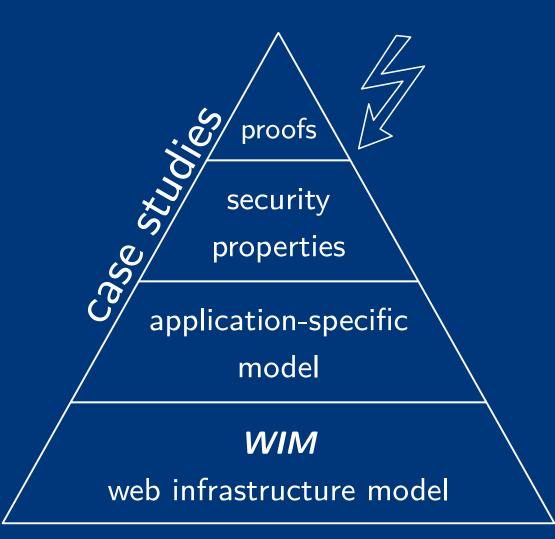
Advantages

This approach can yield...

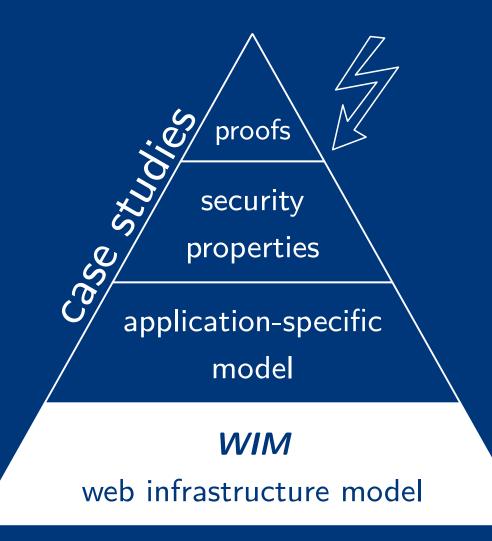
- new attacks and respective fixes
- strong security guarantees excluding even unknown types of attacks



An Expressive Formal Model of the Web Infrastructure



An Expressive Formal Model of the Web Infrastructure



[Kerschbaum 2007]

Analysis of CSRF protection in the Alloy model checker

- [Akhawe, Barth, Lam, Mitchell, Song 2010]
 First formal "web model", in Alloy, five case studies
 - [Bansal, Bhargavan, Maffeis et al. 2012, 2013, 2014] Formal web model with many web features, based on ProVerif tool, new attacks on encrypted cloud storage and OAuth 2.0

Very limited web models

Limitations and constraints of tools (e.g., encoding of messages/terms and data structures)

HTTPTransaction1 \$trSchema: d/HTTP (\$r) HTTPRequest1 HTTPRequest0 HTTPResponse1 headers path HOME0 HOME0 HTTPResponse1 headers path HOME0 HOME1

Our approach: goal was a very detailed, close-to-standards web model, (started with) pen-and-paper.

Further Related Work (Formal Analysis)

- ► [Kumar et al., 2011-2014]: Alloy-based with BAN logic
- ► [Bai et al., 2013]: AuthScan + ProVerif
- ► [Bohannon and Pierce, 2010]
 - "Featherweight Firefox"
 - Information Flow tracking in web browser core
 - No security policies by default
- [Sabelfeld et al. 2016]: Information-flow security for JavaScript and its APIs
- ► [Börger et al., 2012]
 - Abstract State Machines
 - Focus on web server, limited browser model

The Web Infrastructure Model WIM

Detailed, comprehensive, and precise formal model

Network interactions Attacker behavior

DNS servers

Generic web server model

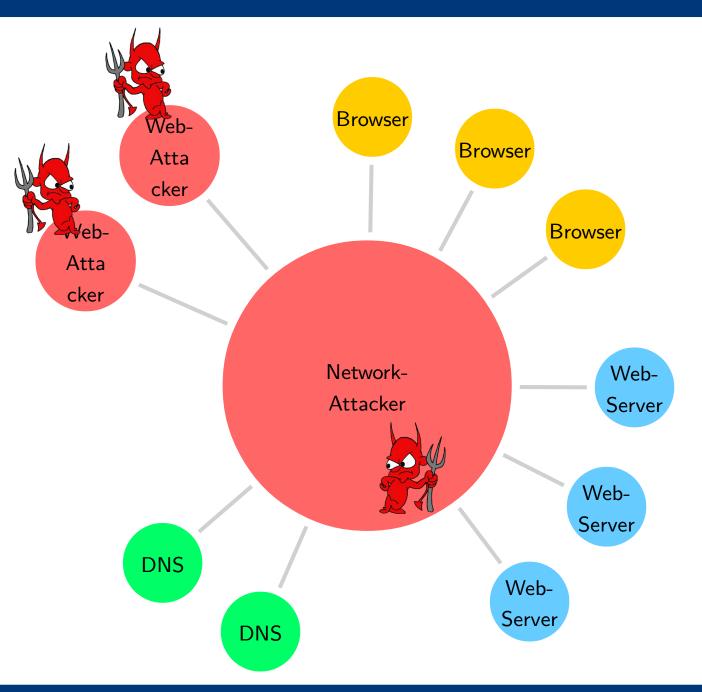
Web browsers

- Summarizes and condenses relevant standards
- Solid basis for security and privacy analyses of web standards and applications

Reference model

developers, researchers, teaching, and tool-based analysis

WIM Network Model and Attackers



Dolev-Yao-Style Model:

- Messages are terms
- Attacker, Browsers, Servers,
 Scripts (honest or malicious)
 are Dolev-Yao processes
- Not just a standard Dolev-Yao model for protocol analysis, but rather covers web features, close to web standards.

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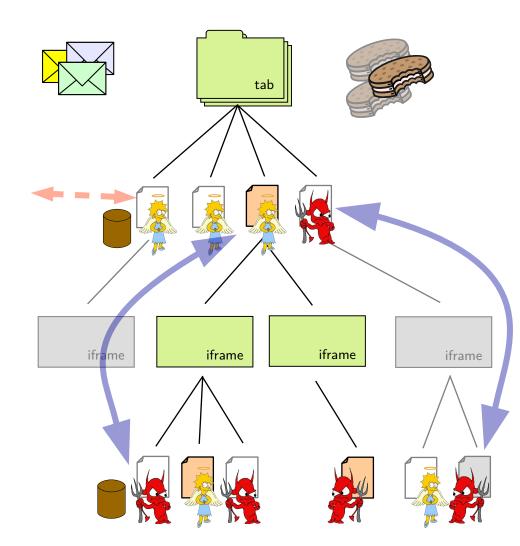
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WIM Web Browser Model



Including ...

• DNS, HTTP, HTTPS



Origin: https://example.com

- window & document structure
- scripts
 (honest and malicious)
- web storage & cookies
- web messaging & XHR
- message headers
 (Origin, STS, Location, Referer, ...)
- redirections
- security policies
- WebRTC
- dynamic corruption





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WIM Web Browser Model - Example

quite complex rules

```
Algorithm 8 Web Browser Model: Process an HTTP response.
```

- 1: function PROCESSRESPONSE(response, reference, request, requestUrl, key, f, s')
- 2: **if** Set-Cookie \in *response*.headers **then**
- 3: for each $c \in \langle \rangle$ response.headers [Set-Cookie], $c \in$ Cookies do
- 4: **let** s'.cookies[request.host]
 - \rightarrow := AddCookie(s'.cookies[request.host], c)
- 5: **if** Strict-Transport-Security \in response.headers \land requestUrl.protocol \equiv S then
- 6: **let** $s'.sts := s'.sts + \langle \rangle$ request.host
- 7: **if** Referer \in *request*.headers **then**

```
8: let referrer := request.headers[Referer]
```

9: **else**

10:

14:

18:

20:

```
let referrer := \bot
```

- 11: **if** Location \in response.headers \land response.status \in {303, 307} **then**
- 12: **let** *url* := *response*.headers [Location]

```
13: if url.fragment \equiv \bot then
```

```
let url.fragment := requestUrl.fragment
```

- 15: **let** *method'* := *request*.method
- 16: **let** body' := request.body
- 17: **if** Origin \in request.headers **then**
 - $let \ origin := \langle request.headers[Origin], \langle request.host, url.protocol \rangle \rangle$

19: else

let $origin := \bot$

- 21: **if** response.status $\equiv 303 \land request.method \notin \{GET, HEAD\}$ then
- 22: **let** method' := GET
- 23: **let** $body' := \langle \rangle$

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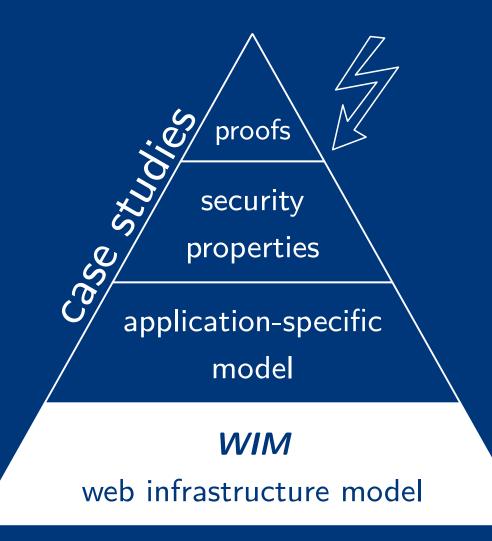
developers, researchers, teaching, and tool-based analysis

Limitations

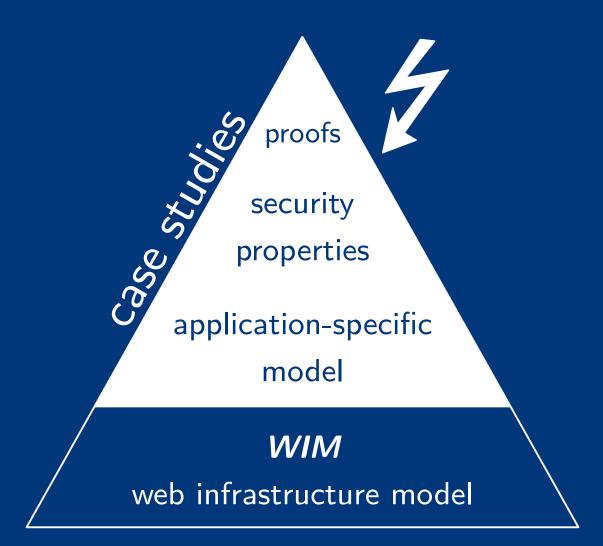
- ► No language details
- No user interface details (e.g., no clickjacking attacks)
- No byte-level attacks (e.g., buffer overflows)
- Abstract view on cryptography and TLS

Model can in principle be extended to capture these aspects as well. Trade-off: comprehensiveness vs. simplicity

An Expressive Formal Model of the Web Infrastructure



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WIM Case Studies

Web single sign-on (SSO) systems

TripAdvisor - Registration - Mozilla Firefox	- + ×	
		Client)
 Where are you going? What are protocol flows Weich and the protocol	Facebook - Mozilla Firefox - + × 3.)facebook.com/login.php?skip_api_login=1&api_k€ Google 8	
	Log In Cancel	

WIM Case Studies



Mozilla BrowserID

SPRESSO https://spresso.me



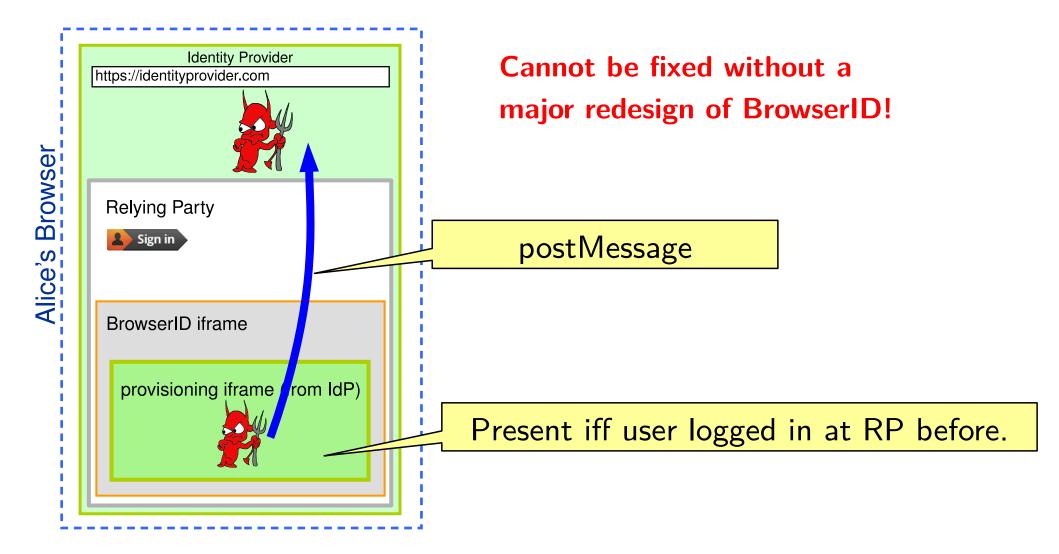
OAuth 2.0



- Discovered severe attacks against authentication
- After fixes: Proof of security
- Special feature privacy: broken beyond repair

BrowserID: Privacy Attack

Information is leaked by the **window structure** in the user's browser:



WIM Case Studies

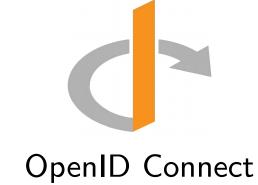


Mozilla BrowserID

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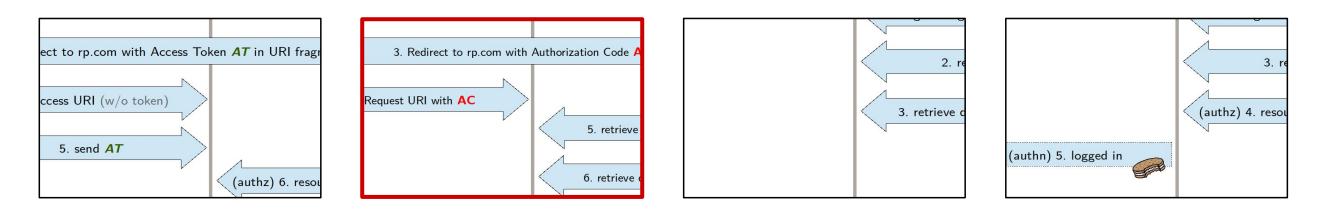
- SSO framework used for authorization/authentication
- Specified by IETF (RFC6749), very widely used

(e.g., F Log in With Facebook

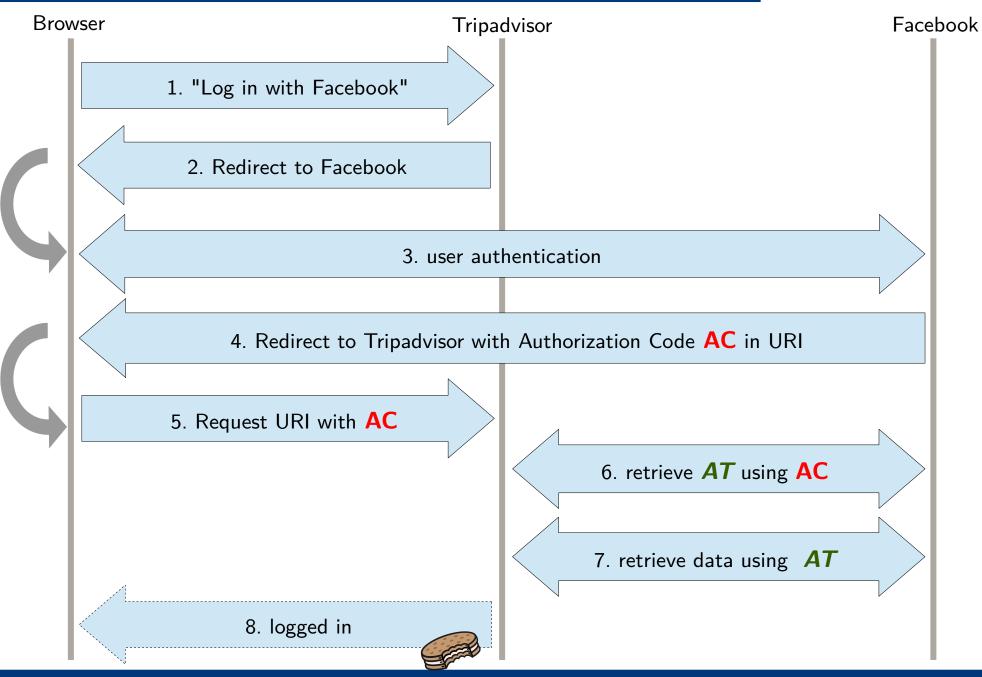
► Many "variables":

optional parameters, *public* and *confidential* clients, etc.

► Four different modes of interaction (*grants*)

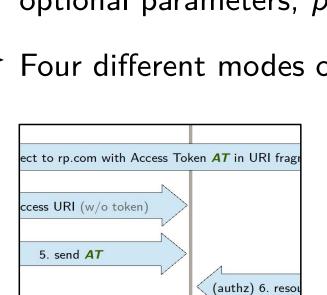


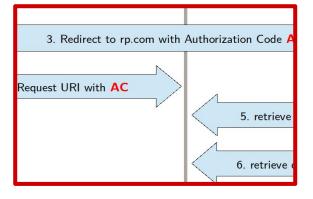
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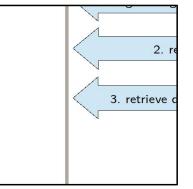


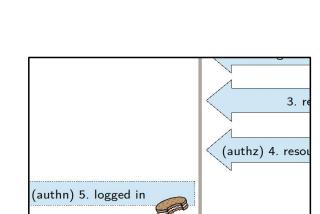
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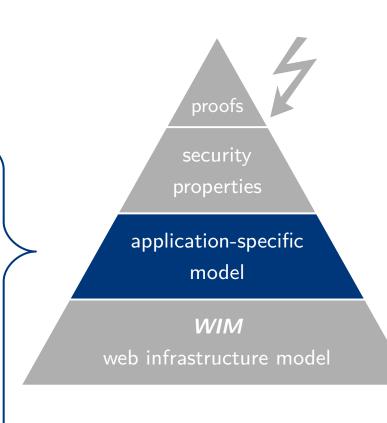
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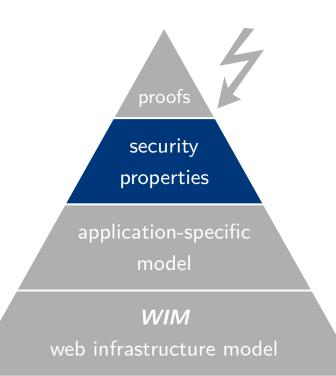






Authentication

Definition 56 (Authentication Property). Let $OAuthWS^n$ be an OAuth web system with a network attacker. We say that $OAuthWS^n$ is secure w.r.t. authentication iff for every run ρ of $OAuthWS^n$, every state (S^j, E^j, N^j) in ρ , every $r \in \text{Clients}$ that is honest in S^j , every $i \in \text{OAP}$, every $g \in \text{dom}(i)$, every $u \in S$, every client service token of the form $\langle n, \langle u, g \rangle \rangle$ recorded in $S^j(r)$.serviceTokens, and n being derivable from the attackers knowledge in S^j (i.e., $n \in d_{\emptyset}(S^j(\text{attacker}))$), then the browser b owning u is fully corrupted in S^j (i.e., the value of *isCorrupted* is FULLCORRUPT), some $r' \in \text{trustedClients}(\text{secretOfID}(\langle u, g \rangle))$ is corrupted in S^j , or i is corrupted in S^j .



Authorization

Definition 55 (Authorization Property). Let $OAuthWS^n$ be an OAuth web system with a network attacker. We say that $OAuthWS^n$ is secure w.r.t. authorization iff for every run ρ of $OAuthWS^n$, every state (S^j, E^j, N^j) in ρ , every OAP $i \in OAP$, every $r \in Clients \cup \{\bot\}$ with rbeing honest in S^j unless $r = \bot$, every $u \in ID \cup \{\bot\}$, for n = resourceOf(i, r, u), n is derivable from the attackers knowledge in S^j (i.e., $n \in d_{\emptyset}(S^j(attacker)))$, it follows that

- 1. *i* is corrupted in S^j , or
- 2. $u \neq \perp$ and (i) the browser b owning u is fully corrupted in S^j or (ii) some $r' \in$ trustedClients(secretOfID(u)) is corrupted in S^j .

OAuth 2.0: Security Properties

Session Integrity for authentication

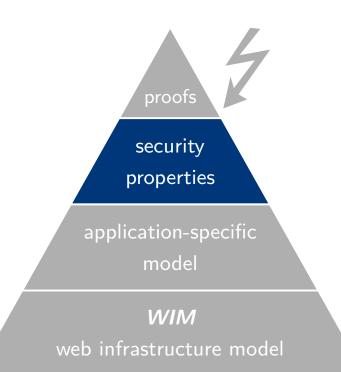
Definition 64 (Session Integrity for Authentication). Let $OAuthWS^w$ be an OAuth web system with web attackers. We say that $OAuthWS^w$ is secure w.r.t. session integrity for authentication iff for every run ρ of $OAuthWS^w$, every processing step Q_{login} in ρ , every browser b that is honest in Q_{login} , every $r \in \text{Clients}$ that is honest in Q_{login} , every $i \in \text{OAP}$, every identity $\langle u, g \rangle$, the following holds true: If in Q_{login} a service token of the form $\langle n, \langle \langle u', g' \rangle, m \rangle \rangle$ for a domain $m \in \text{dom}(i)$ and some n, u', g' is created in r (in Line 38 of Algorithm B.4) and n is sent to the browser b, then

```
(a) there is an OAuth Session o \in \mathsf{OASessions}(\rho, b, r, i), and
```

(b) if i is honest in Q_{login} then Q_{login} is in o and we have that

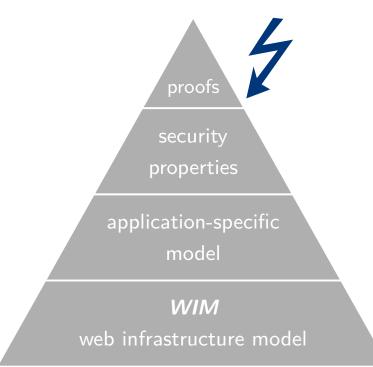
 $\left(\mathsf{selected}_{\mathrm{ia}}(o, b, r, \langle u, g \rangle) \lor \mathsf{selected}_{\mathrm{nia}}(o, b, r, \langle u, g \rangle)\right) \iff \left(\langle u, g \rangle \equiv \langle u', g' \rangle\right).$

 Session Integrity for authorization (similar to above)



OAuth 2.0: New Attacks

OAuth 2.0 had been analyzed many times before, but not in a comprehensive formal model.



Further Related Work (OAuth 2.0)

- ► [Bansal et al., 2012-2014]
- ► [Wang et al., 2013]
 - "Explicating SDKs"
 - Boogie/Corral
 - Extraction of SDK logic, definition of security properties, addition of assume statements, code verification.
- ► [Chari, Jutla, Roy, 2011]
 - UC model analysis of OAuth Authorization Code Grant
 - No web features
- Several empirical studies, focussed on typical implementation errors

Further Related Work (OpenID Connect)

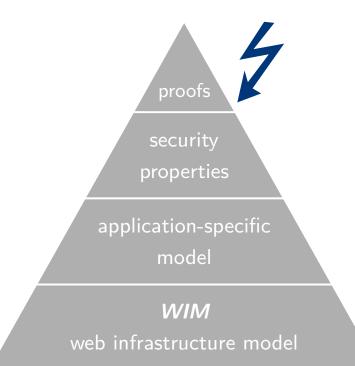
- ► [Mladenov et al., 2016]
 - Specific variant of the IDP Mix-Up attack
 - No formal model
- ► [Li, Mitchell, 2016]
 - Implementation errors in deployments of Google Sign-In

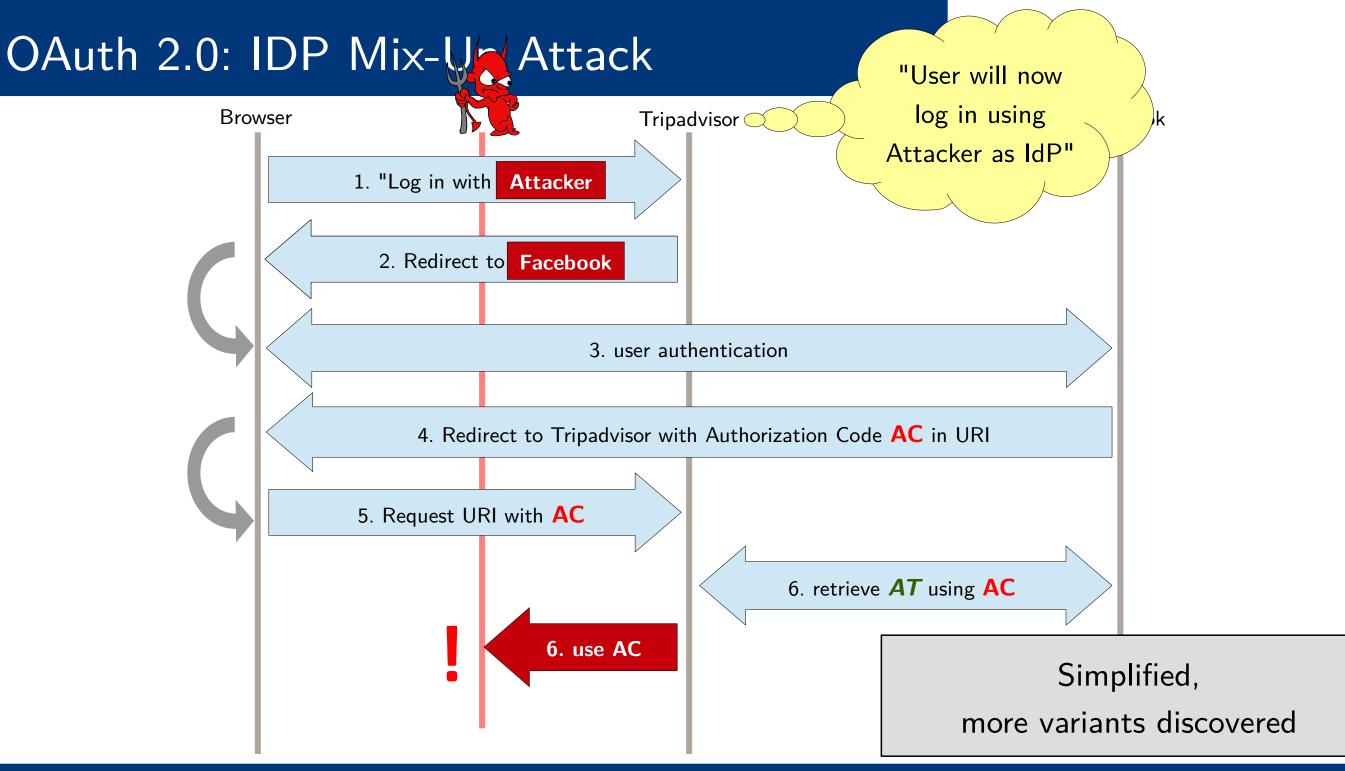
OAuth 2.0: New Attacks

OAuth 2.0 had been analyzed many times before, but not in a comprehensive formal model.

New attacks:

- ► 307 Redirect Attack
- Identity Provider Mix-Up Attack (new class of attacks)
- State Leak Attack
- Naïve Client Session Integrity Attack
- Across Identity Provider State Reuse Attack



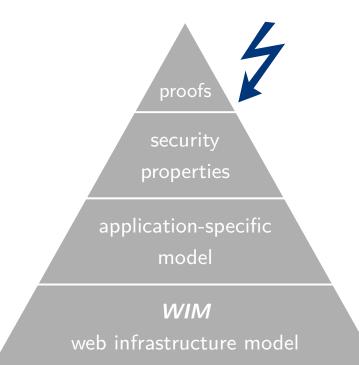


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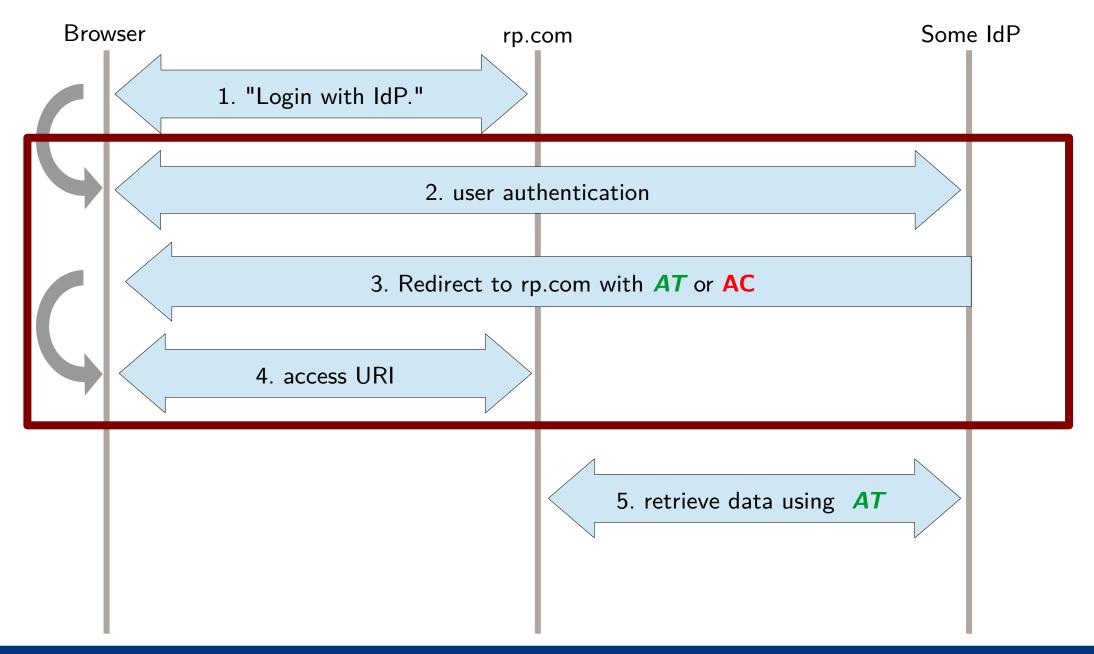
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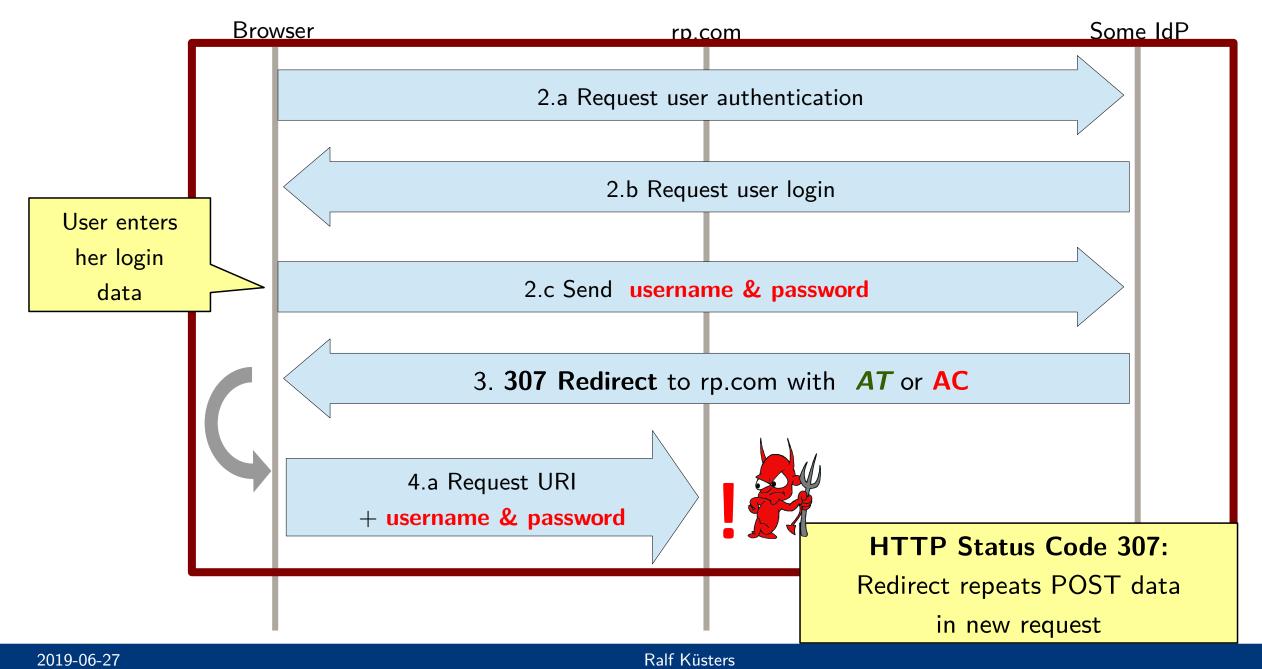
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OAuth: 307 Redirect Attack (I)



OAuth: 307 Redirect Attack (II)



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OAuth 2.0: Proof of Security

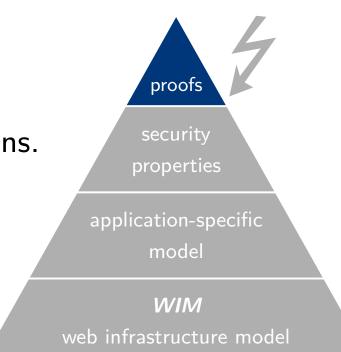
Proof based on our model of OAuth 2.0 with all grant types and options.

Assumptions:

- Adherence to web best practices (e.g., regarding session handling)
- Adoption of our implementation guidelines

 (e.g., no 3rd party scripts on certain web pages)
- Fixes against previously known and new attacks

Theorem 1. Let $OAuthWS^n$ be an OAuth web system with a network attacker, then $OAuthWS^n$ is secure w.r.t. authorization and secure w.r.t. authentication. Let $OAuthWS^w$ be an OAuth web system with web attackers, then $OAuthWS^w$ is secure w.r.t. session integrity for authorization and authentication.



OAuth 2.0: Impact

- Disclosed OAuth attacks to the IETF Web Authorization Working Group in late 2015
- Emergency meeting with the working group four weeks later
- Initiated the OAuth Security Workshop (OSW) to foster the exchange between researchers, standardization groups, and industry
- Joined the working group to codify the fixes into a new RFC: OAuth 2.0 Security Best Current Practice [draft-ietf-oauth-security-topics]

WIM Case Studies



Mozilla BrowserID

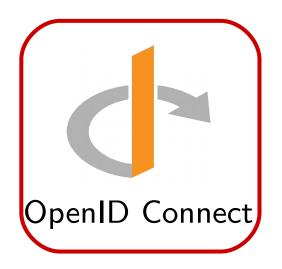
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- After fixes: Proof of authentication
- Special feature privacy: broken beyond repair



- Designed from scratch
- First formalized in
 WIM, then
 - implemented
- First SSO with proven privacy and security



- Found several new attacks
- Developed fixes and implementation guidelines
- Proof of security



OpenID Connect

- OAuth 2.0 was built for authorization, not authentication
- OpenID Connect: "Identity Layer" for OAuth 2.0 to solve this
- Includes new extensions:
 - Automatic discovery of identity providers
 - Dynamic registration of clients at identity providers
- New token type ("id token")
- Cryptographic mechanisms, e.g., signed id token

Out of scope of plain OAuth 2.0

Results:

- All newly discovered OAuth attacks apply to OpenID Connect as well
- Implementation guidelines to avoid known attacks
- Proof of security (authentication, authorization, session integrity) including discovery and dynamic registration extensions

Theorem 2 (Security of OpenID Connect). Let $OIDCWS^n$ be an OIDC web system with a network attacker. Then, $OIDCWS^n$ is secure w.r.t. authentication and authorization. Let $OIDCWS^w$ be an OIDC web system with web attackers. Then, $OIDCWS^w$ is secure w.r.t. session integrity for authentication and authorization.

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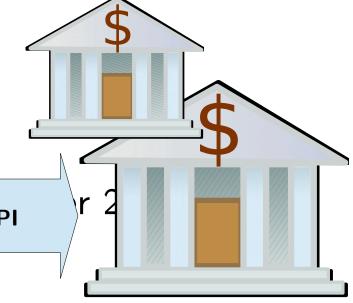
- Including extensions
- Developed best practices against known attacks
- Proof of security

Most recent case study: **Financial-grade API (FAPI)**

- Authorization and authentication in high-risk scenarios
- Laws and activities for opening financial services to third-party providers
 - <u>OpenBanking UK: Financial-grade API already</u> ted by major banks in the Authorize app: **Financial-grade API** Access to banking account

other countries follow sin

proaches

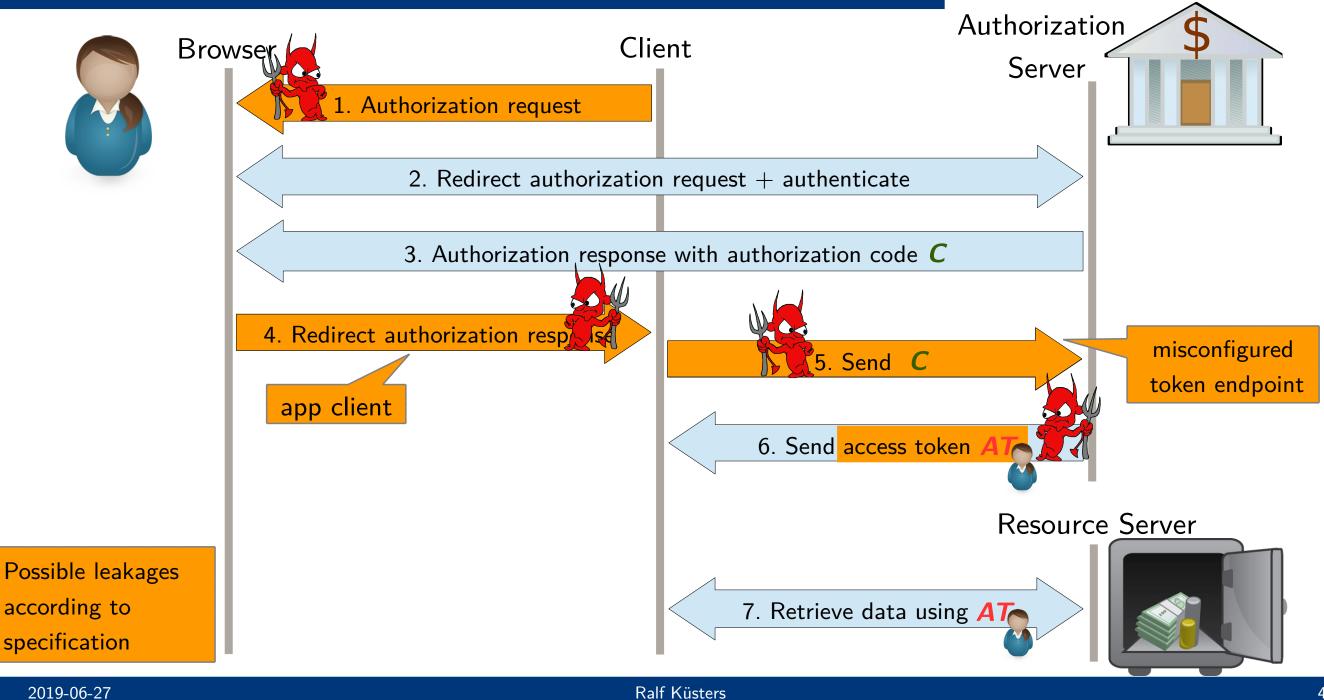


Overview FAPI

OpenID Financial-grade API:

- Hardened version of **OAuth 2.0** for high-risk use-cases
- New mechanisms: OAuth 2.0 Token Binding, OAuth 2.0 Mutual TLS, Proof Key for Code Exchange, JWT Secured Authorization Response Mode

FAPI: Attacker Model



Overview FAPI

- OpenID Financial-grade API:
 - Hardened version of **OAuth 2.0** for high-risk use-cases
 - New mechanisms: OAuth 2.0 Token Binding, OAuth 2.0 Mutual TLS, Proof Key for Code Exchange, JWT Secured Authorization Response Mode
- Our Work: formal security analysis of the Financial-grade API
 - Formal model of the Financial-grade API based on the **Web Infrastructure Model**
 - Precise definition of **security properties**
 - During formal analysis: **found several attacks** bypassing the new mechanisms
 - **Proof of security** for the fixed Financial-grade API
- Collaborating with OpenID Foundation to fix the standard

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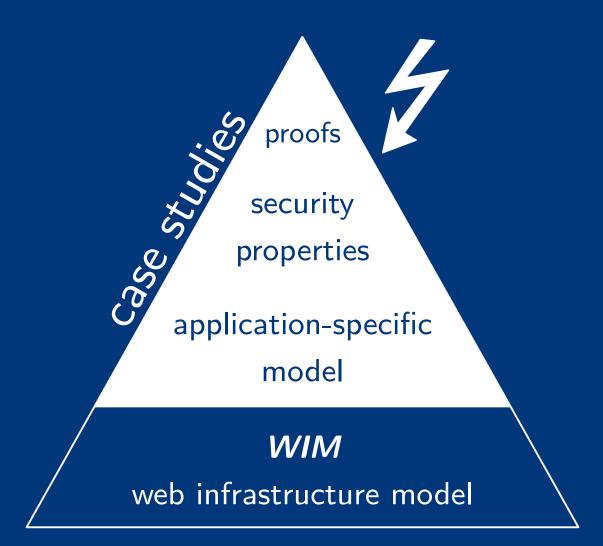


OpenID Connect

- Including extensions
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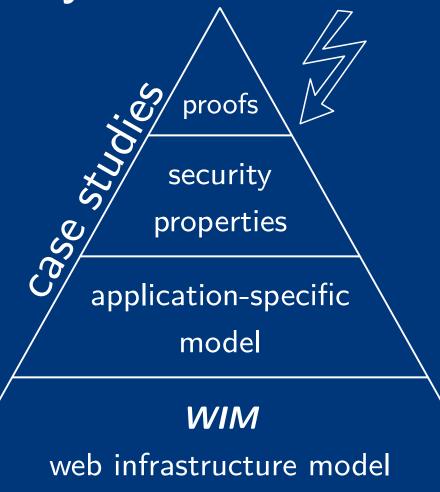
Most recent case study: Financial-grade API (FAPI)

An Expressive Formal Model of the Web Infrastructure



WIM: An Expressive Formal Model of the Web Infrastructure

Thank you!



Most detailed and comprehensive formal model of the web infrastructure so far

- Case studies with real-world impact
- New classes of attacks
- Formal proofs of web security with very high level of detail
- Designed first privacy-preserving SSO system: SPRESSO
- Currently: mechanized model, in collaboration with Bhargavan et al.